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L17

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L6: Entry 1 of 4

File: USPT

Jan 15, 2002

DOCUMENT-IDENTIFIER: US 6339837 B1

TITLE: Hybrid method for design verification

Detailed Description Text (21):

FIG. 2 illustrates a verification engine in the first embodiment, comprising four steps: an interpreter 201, a cycle counter 202, an unroller 203, and an equivalence checker 204. Interpreter 201 translates verification structure 111 from a hardware description language into a data structure representing a network of gates and registers. Cycle counter 202 determines the clock cycle number involved in the behavior being simulated up to the ending indicator. Unroller 203 traverses the network of gates and registers in the result of interpreter 201 from result flag 104, with the clock cycle number from cycle counter 202, to free variables 101 and builds a data structure representing a network of gates. Unroller 203 copies the contributing gates for each clock cycle and connects the consecutive copies by replacing each register with a wire between its data input and its data output. For the initial cycle, each register's data output takes the register's initial value if there is any, or otherwise one of free variables 101. For the final cycle, the register inputs are ignored if they are not connected to result flag 104. Equivalence checker 204 determines whether the network of gates makes every copy of result flag 104 generated by unroller 203 equivalent to the predetermined value of result flag 104. The answer from equivalence checker 204 is taken to a decision point 205, which declares the verification a success 206 if it is all positive. Otherwise it declares the verification a failure 207. All steps before taking the answer from equivalence checker 204 can include Boolean simplification techniques for the sake of reducing the amount of computation.

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S, 146/460

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L4: Entry 2 of 14

File: USPT

Nov 30, 2004

DOCUMENT-IDENTIFIER: US 6826518 B1

TITLE: Method for distributed agent-based non-expert simulation of manufacturing process behavior

Brief Summary Text (21):

From the agent perspective, a method for distributed agent-based simulation of manufacturing process behavior, the simulation having a plurality of agents corresponding to individual processes forming a manufacturing technique, comprises the steps of: receiving a message from an agent; identifying in the received message a discrete event selected from the group consisting of a clock tick event, a resources received event, and a request for output production event; causing an associated process to perform an activity in response to the identified event; and, messaging an adjacent agent in response to the identified event.

CLAIMS:

7. A method for distributed agent-based simulation of manufacturing process behavior, the simulation having a plurality of agents corresponding to individual processes forming a manufacturing technique, the method comprising the steps of: receiving a message from an agent associated with one of a plurality of different manufacturing techniques comprising a pull, a push, or a takt manufacturing technique, wherein the agent is responsive to a discrete event selected from the group consisting of a clock tick message, a resources received message, and a request for output production message; identifying in said received message a discrete event selected from the group consisting of a clock tick event, a resources received event, and a request for output production event; causing an associated process to perform an activity in response to said identified event; and, messaging an adjacent agent in response to said identified event.

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L17: Entry 9 of 20

File: USPT

Apr 18, 2000

DOCUMENT-IDENTIFIER: US 6052524 A

TITLE: System and method for simulation of integrated hardware and software components

Brief Summary Text (14):

There have been attempts to overcome some of the timing problems associated with simulators and emulators using hardware verification languages including: "Logic Simulation Using a Hardware Accelerator Together with an Automated Error Event Isolation and Trace Facility," U.S. Pat. No. 5,146,460, by Ackerman et al; "Timing Analysis for Logic Optimization Using Target Library Delay Value," U.S. Pat. No. 5,475,605, by Lin; and "Hardware-Software Debugger Using Simulation Speed Enhancing Techniques Including Skipping Unnecessary Bus Cycles, Avoiding Instruction Fetch Simulation, Eliminating the Need for Explicit Clock Pulse Generation and Caching Results of Instruction Decoding," U.S. Pat. No. 5,678,028, by Bershteyn, et al. However, as electronic devices continually increase in speed and require different types of high speed and low speed memory, such attempts to improve the timing problems still do not provide adequate simulation of timing interactions between hardware and software components in an electronic device.

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L17: Entry 18 of 20

File: USPT

Sep 8, 1992

DOCUMENT-IDENTIFIER: US 5146460 A

TITLE: Logic simulation using a hardware accelerator together with an automated error event isolation and trace facility

Abstract Text (1):

Software simulators of logic design circuits run slowly but are capable of providing very finely detailed error trace analyses. On the other hand, hardware accelerators operating to perform similar functions are very fast in their execution but are not capable of practically isolating error states or other critical conditions. Accordingly, the present invention provides an interactive system combining software simulators and hardware accelerators so that when desired test results do not favorably compare with simulated results, a mechanism is provided for storing the current hardware accelerator state and restoring the accelerator to a previous checkpoint state which has been saved as a result of a prior periodic interruption. The hardware accelerator is then operated for a time sufficient to bring it up to a state that occurs just before the detected miscomparison. At this point, state information from the hardware accelerator is supplied to a software simulator for detailed error analysis and fault tracing. The hardware accelerator may then be restarted where it left off or with a different task. In this way, optimal utilization is made of expensive hardware accelerator resources while nonetheless providing error event isolation and the ability to generate detailed traces of the simulated behavior.

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L8: Entry 1 of 1

File: USPT

Aug 22, 2000

DOCUMENT-IDENTIFIER: US 6106298 A

TITLE: Reconfigurable easily deployable simulator

Abstract Text (1):

A modular, reconfigurable and easily deployable simulator system includes a plurality of vehicle personality kits and at least one frame to support a single crewstation. Members of a selected kit are coupled to the frame. A multi-platform control system provides on-the-window and in vehicle instrument displays in real-time in response to the crewperson manipulating platform controls. An editor with a graphical user interface provides a capability to modify program segments of the control system thereby altering the behavior of the simulated platform.

Detailed Description Text (75):

The configuration file specifies the segments and interconnecting messages which ultimately carry out the control process leading to behavior which simulates the selected platform. A configuration file includes a manager segment. The manager segment issues a load configuration call. The configuration file is then read in the segment executable so called. The called segments in turn issue enter configuration calls to register with the CSL 16. They then initiate various generic procedure calls to access message data. When the segments are all initialized, the simulation process can begin.

Detailed Description Text (85):

A particular processor can be specified to which a given segment will be associated. Absent binding a segment to a processor, CSL 16 and the operating system may move a particular segment to a different host for execution. A path parameter specifies a UNIX path for the segment executable or script file that will run the segment executable. A priority parameter controls the priority that the segment will receive relative to other processors or segments being run on the host. A lock parameter makes it possible to lock a segment's text and data area thereby making it immune to routine swapping by the operating system. A list of message names must be provided along with the type of access to indicate those messages that the segment needs access to. Finally, X and Y position parameters for the configuration editor graphical user interface must also be specified.

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File: USPT

Mar 27, 2001

DOCUMENT-IDENTIFIER: US 6208955 B1

TITLE: Distributed maintenance system based on causal networks

Detailed Description Text (35):

In step 610 detailed simulation for model testing and validation is performed. The causal network diagnostic model operates on a static-input-variable-state snapshot to produce a diagnostic output. This process can be performed either through a manual mode or by executing simulation script files that specify the timing constraints for individual modeling components. In the manual mode, failure sources can be selected to change their states so as to simulate the rippling effect of a failure source through avionics system 100. Thus, system level evaluation for multiple system states and diagnostic outputs is possible in a short period of time. For the simulation script method, an input/output (I/O) shell simulator is used to simulate inputs into the model domain and control the timing relationships. It is the simulator shell's job to activate the diagnostic process when an appropriate system state has been reached. Activation of the diagnostics is therefore event-driven.

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